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  - ▶ expanding frontier – “greedy” – propagate a wave 1 edge-distance at a time.

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**Enqueue(Q,v)**: insert element  $v$  into the queue  $Q$   
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- ▶ There are several efficient ways to implement queues and stacks, see section 10.1.

## Breadth-First Search (BFS)

```
BFS(G,s)
for each vertex u in V-{s}
    d[u] = +infty
endfor
d[s] = 0
Q = empty // create FIFO queue
Enqueue(Q, s)
while Q not empty
    u = Dequeue(Q)
    for each v in Adj[u]
        if d[v] = +infty
            d[v] = d[u] + 1
            Enqueue(Q, v)
        endif
    endfor
endwhile
return d
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- ▶ Correctness of BFS  
shortest path proof – see pp.597-600 of [CLRS,3rd ed.]  
similar with weighted edges – Dijkstra's algorithm – *to be discussed*