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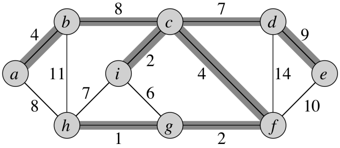
  Example



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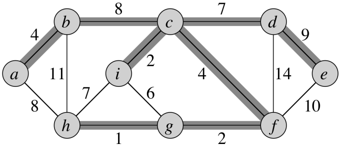

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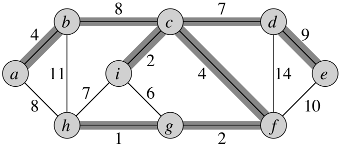

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Example: $w(T) = 37$.

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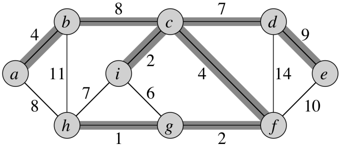

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- MST is not necessarily unique.
  For simplicity in theory, assume all edge weight distinct, and therefore, has a unique MST.

MST

Basic idea of computing ("growing") a MST:

- construct the MST by successively select edges to include in the tree

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One of the most famous greedy algorithms, along with Huffman coding

MST

Two basic properties:

1. Optimal substructure: optimal tree contains optimal subtrees.

[1] The subgraph G_1 is induced by vertices in T_1, i.e., $V_1 = \{\text{vertices in } T_1\}$ and $E_1 = \{(x, y) \in E; x, y \in V_1\}$. Similarly for G_2.

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   Let T be a MST of $G = (V, E)$. Removing (u, v) of T partitions T into two trees T_1 and T_2. Then T_1 is a MST of $G_1 = (V_1, E_1)$ and T_2 is a MST of $G_2 = (V_2, E_2)$.[1]

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   Proof. Note that

   $$w(T) = w(T_1) + w(u, v) + w(T_2).$$

   There cannot be a better subtree than T_1 or T_2, otherwise T would be suboptimal.

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[2] Note: there is an abuse of notation here that we will view A as being both edges and vertices.

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   Let T be a MST of $G = (V, E)$, $A \subseteq T$ be a subtree of T, and (u, v) be min-weight edge in G connecting A and $V - A$. Then $(u, v) \in T$.[2]

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   Proof. If $(u, v) \notin T$, then

   - $(u, v) \cup T$ forms a cycle,
   - replace one of edges of T by (u, v) form a new tree T
   - this is contradiction to T is MST

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MST

Prim's algorithm

- Basic idea:
  - starts from an arbitrary root r

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- How to find the next lightest edge quickly?

  Answer: use a **priority queue**

Review: Priority Queue

A priority queue maintains a set S of elements, each with an associated value called a "key", and supports the following operations:

- Search(S,k):
  returns x in S with key[x] = k
- Insert(S, x)/Delete(S, x):
  inserts/deletes the element x into the set S
- Maximum(S)/Minimum(S):
  returns x in S with largest/smallest key
- Extract-max(S)/Extract-min(S):
  removes and returns x in S with largest/smallest key
- Increase-key(S, x, k)/Decrease-key(S, x, k):
  increases/decreases the value of element x's key to the new value k

Recall that the priority queue has been used in Huffman coding.

MST

```
MST-Prim(G, w, r)
Q = empty
for each vertex u in V
    key[u] = infty  // min. weight of any edge (w,u) and w in A
    pi[u] = nil     // parent of u
    Insert(Q, u)
endfor
Decrease-key(Q,r,0)
while Q not empty
    u = Extract-Min(Q)
    for each v in Adj[u]
        if (v in Q) and (w(u,v) < key[v])
            Decrease-key(Q, v, w(u,v))
            pi[v] = u   // parent of v
       endif
    endfor
endwhile
return A = { (v, pi[v]): v in V-{r} }   // MST
```

MST

Run and *illustrate* Prim's algorithm

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       endif
    endfor
endwhile
return A = { (v, pi[v]): v in V-{r} }   // MST
```

MST

Prim's algorithm

1. Run and *illustrate* Prim's algorithm
2. Running time:
   - depends on how the priorty queue Q is implemented
   - Suppose Q is a binary heap (see Section 6.1)
     - Initialize Q and the first for loop: $O(|V| \lg |V|)$
     - Decrease key of root r: $O(\lg |V|)$
     - While-loop:
       a) $|V|$ Extract-Min calls: $O(|V| \lg |V|)$
       b) $\leq |E|$ Decrease-Key calls: $O(|E| \lg |E|)$
   - Total: $O(|E| \lg |V|)$
   - *Note: G is connected,* $\lg |E| = \Theta(\lg |V|)$

MST

Kruskal's algorithm

- Basic idea:
  - scan edges in increasing of weight
  - put edge in if no loop created

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- Why does this result in MST?
  Answer: min-weight edge is always in MST (the greedy-choice property).

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- Basic idea:
  - scan edges in increasing of weight
  - put edge in if no loop created
- Why does this result in MST?
  Answer: min-weight edge is always in MST (the greedy-choice property).
- How to make sure "no loop created"?
  use "disjoint-set" data structure

Review: Disjoint-Set

Disjoint-Set maintains a collection of $S = \{S_1, S_2, ...S_k\}$ of disjoint dynamic sets. Each set is identified by a representative, which is some member of the set.

A disjoint-set data structure supports the following operations:

- Make-set(x):
  creates a new set whose only member (and thus representative) is x.
- Union(x, y):
  unites the sets that contain x and y, say S_x and S_y, into a new set that is the union of these two sets: $S_x \cup S_y$. The representative is any member of $S_x \cup S_y$.
- Find-set(x):
  returns (a pointer to) the representative of the (unique) set containing x.

To learn more about the disjoint-set data structure, see Chapter 21.

MST

```
MST-Kruskal(G, w)
A = emtpy
for each vertex v in V
    Make-set(v)
endfor
Sort the edges E in nondecreasing order by w
for each edge (u,v) in E, taken in nondecreasing order by w
    if Find-set(u) \= Find-set(v)
        A = A U {(u,v)}
        Union(u,v)
    endif
endfor
return A
```

MST

Run and *illustrate* Prim's algorithm

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```

MST

Kruskal's algorithm

1. Run and *illustrate* Prim's algorithm
2. Running time:
    - depends on the implementation of the disjoint-set
    - Sort: $\Theta(|E| \lg |E|)$
    - $|V|$ Make-Set ops
    - $2|E|$ Find-Set ops
    - $|V| - 1$ Union ops
    - Total: $O(|E| \lg |V|)$
    - *Note: G is connected,* $\lg |E| = \Theta(\lg |V|)$