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- Edit distance is the minimum number of edits - insertions, deletions and substitutions of characters - need to transform the first string into the second. e.g. a spell checker.


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- Subproblem:
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- How to express $e(i, j)$ in terms of its subproblems, recursively?
- key observation: the rightmost column of an alignment of $x[1 \cdots i]$ and $y[1 \cdots j]$ can only be one of the following three cases:
Case 1
Case 2
Case 3
$x[i]$
- 

or
-
$y[j]$
or
$x[i]$
$y[j]$

## Edit distance

- By the above key observation, then

$$
e(i, j)=\min \{\underbrace{1+e(i-1, j)}_{\text {case } 1}, \underbrace{1+e(i, j-1)}_{\text {case } 2}, \underbrace{\operatorname{diff}(i, j)+e(i-1, j-1)}_{\text {case } 3}\}
$$

where

$$
\operatorname{diff}(i, j)= \begin{cases}0 & \text { if } x[i]=y[j] \\ 1 & \text { if } x[i] \neq y[j]\end{cases}
$$

- Question: how to find the corresponding optimal alignment?


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|  |  | s | u | n | n | y |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | 0 | 1 | 2 | 3 | 4 | 5 |
| s | 1 | 0 | 1 | 2 | 3 | 4 |
| n | 2 | 1 | 1 | 1 | 2 | 3 |
| o | 3 | 2 | 2 | 2 | 2 | 3 |
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| o | 3 | 2 | 2 | 2 | 2 | 3 |
| w | 4 | 3 | 3 | 3 | 3 | 3 |
| y | 5 | 4 | 4 | 4 | 4 | 3 |

Therefore, the edit distance between $x$ and $y=e(5,5)=3$.

## Edit distance

Example 2. $x=$ 'heroically', $y=$ 'scholarly'

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|  |  | s | c | h | o | I | a | r | I | y |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 |
| h | 1 | 1 | 2 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
| e | 2 | 2 | 2 | 3 | 3 | 4 | 5 | 6 | 7 | 8 |
| r | 3 | 3 | 3 | 3 | 4 | 4 | 5 | 5 | 6 | 7 |
| o | 4 | 4 | 4 | 4 | 3 | 4 | 5 | 6 | 6 | 7 |
| i | 5 | 5 | 5 | 5 | 4 | 4 | 5 | 6 | 7 | 7 |
| c | 6 | 6 | 5 | 6 | 5 | 5 | 5 | 6 | 7 | 8 |
| a | 7 | 7 | 6 | 6 | 6 | 6 | 5 | 6 | 7 | 8 |
| l | 8 | 8 | 7 | 7 | 7 | 6 | 6 | 6 | 6 | 7 |
| I | 9 | 9 | 8 | 8 | 8 | 7 | 7 | 7 | 6 | 7 |
| y | 10 | 10 | 9 | 9 | 9 | 8 | 8 | 8 | 7 | 6 |

## Edit distance

Example 2. $x=$ 'heroically', $y=$ 'scholarly'

|  |  | s | c | h | o | I | a | r | l | y |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 |
| h | 1 | 1 | 2 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
| e | 2 | 2 | 2 | 3 | 3 | 4 | 5 | 6 | 7 | 8 |
| r | 3 | 3 | 3 | 3 | 4 | 4 | 5 | 5 | 6 | 7 |
| o | 4 | 4 | 4 | 4 | 3 | 4 | 5 | 6 | 6 | 7 |
| i | 5 | 5 | 5 | 5 | 4 | 4 | 5 | 6 | 7 | 7 |
| c | 6 | 6 | 5 | 6 | 5 | 5 | 5 | 6 | 7 | 8 |
| a | 7 | 7 | 6 | 6 | 6 | 6 | 5 | 6 | 7 | 8 |
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Therefore, the edit distance between $x$ and $y=e(10,9)=6$

## Edit distance

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|  | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 |
| h | 1 | 1 | 2 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
| e | 2 | 2 | 2 | 3 | 3 | 4 | 5 | 6 | 7 | 8 |
| r | 3 | 3 | 3 | 3 | 4 | 4 | 5 | 5 | 6 | 7 |
| o | 4 | 4 | 4 | 4 | 3 | 4 | 5 | 6 | 6 | 7 |
| i | 5 | 5 | 5 | 5 | 4 | 4 | 5 | 6 | 7 | 7 |
| c | 6 | 6 | 5 | 6 | 5 | 5 | 5 | 6 | 7 | 8 |
| a | 7 | 7 | 6 | 6 | 6 | 6 | 5 | 6 | 7 | 8 |
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Therefore, the edit distance between $x$ and $y=e(10,9)=6$
Note: $\operatorname{LCS}(x, y)=5$

