## Problem Set $4-$ Due Tuesday, April 27, 2010, at 4:15 pm

Problem 1. Describe a decision procedure to solve the following problem: given a regular expression $\alpha$, is $\alpha$ a shortest regular expression for $L(\alpha)$ ? How efficient is your procedure?

Problem 2. Are the following statements true or false? Either prove the statement or give a counterexample to it.

Part 2A. If $L \cup L^{\prime}$ is regular than $L$ and $L^{\prime}$ are regular.
Part 2B. If $L^{*}$ is regular than $L$ is regular.
Part 2C. If $L L^{\prime}$ is regular than $L$ and $L^{\prime}$ are regular.
Part 2D. If $L$ and $L^{\prime}$ agree on all but a finite number of strings, then one is regular iff the other is regular.

Part 2E. If $R$ is regular, $L$ is not regular, and $L$ and $R$ are disjoint, then $L \cup R$ is not regular.

Problem 3. Define $A=\left\{x \in\{a, b, \not \forall\}^{*}: x\right.$ contains an equal number of $a$ 's and $b$ 's or $x$ contains consecutive $\sharp \mathrm{s}$ or consecutive letters $\}$.

Part 3A. Can you use the pumping lemma to prove that $L$ is not regular? Explain.
Part 3B. Prove that $A$ is not regular. Hint: consider closure under homomorphisms and problem 2E.

Problem 4. Give a context free grammar for $L=\left\{a^{n} b^{m}: n \neq 2 m\right\}$. Try to make your grammar unambiguous - and explain why it is unambiguous.

Problem 5. A regular grammar is a context-free grammar $G=(V, \Sigma, R, S)$ in which every rule is of the form $A \rightarrow \varepsilon$ or $A \rightarrow a B$, where $a$ is a terminal and $A$ and $B$ are variables. Show that $L$ is regular iff $L$ is generated by a regular grammar.

Problem 6. Consider the grammar $G$ defined by $S \rightarrow A A, A \rightarrow A A A|b A| A b \mid a$.
(a) Carefully and precisely describe the $L(G)$ in an easy-to-recognize form.
(b) Is $L(G)$ regular? Prove your answer either way.
(c) Is $G$ ambiguous? Prove your answer either way.
(d) Is $L(G)$ inherently ambiguous? Give a convincing argument either way.

