# Shape Analysis with the Delaunay Triangulation 

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## Shape of a Point Set



## Surface Reconstruction



Input: Samples from object surface.


Output: Polygonal model.

## Point Set Capture



Point Grey Bumblebee

Cyberware model 15


## Applications

> Levoy et al, Stanford

Delson et al, AMNH


Allen, Curless, Popovic, U Wash.

## Power Diagram

Weighted Voronoi diagram. Input: balls.
$\operatorname{Dist}(x$, ball $)=\operatorname{dist}^{2}(x$, center $)$-radius ${ }^{2}$


Polyhedral cells, same algorithm as regular Voronoi diagram (lift to convex hull)

## Alpha-shapes



Overlay Voronoi diagram of balls on union of input balls. Discard features outside of the union.

## Alpha-shapes



Alpha shape is the set of weighted Delaunay features dual to the weighted Voronoi features intersecting union of balls.

## Alpha-shapes

Edelsbrunner, Kirkpatrick, Seidel, 83
Edelsbrunner, 93: Alpha shape is homotopy equivalent to union of balls, close correspondence with union structure.
Edelsbrunner \& Muecke, 94: 3D surface reconstruction.

## Alpha-shape reconstruction



Put small ball around each sample, compute alpha-shape.

## Difficulty



> Usually no ideal choice of radius.

## Ball-pivoting

Bernardini et al, IBM


Fixed-radius ball "rolling" over points selects subset of alpha-shape.

## Medial Axis

Medial axis is set of points with more than one closest surface point.

Blum, 67


## 3D Medial Axis

Medial axis of a surface forms a dual surface.


## Medial Axis



Maximal ball avoiding surface is a medial ball.

Every solid is a union of balls!

## Relation to Voronoi



Voronoi balls approximate medial balls.

For dense surface samples in 2D, all Voronoi vertices lie near medial axis.
Ogniewicz, 92

## Convergence



> In 2D, set of Voronoi vertices converges to the medial axis as sampling density increases.

## Discrete unions of balls



Voronoi balls approximate the object and its complement.

## 2D Curve Reconstruction



Blue Delaunay edges reconstruct the curve, pink triangulate interior/exterior.

Many algorithms, with proofs, for coloring edges.

## 2D Medial Reconstruction



Pink approximate medial axis.

By nerve theorem, approximation is homotopy equivalent to object and its complement.

## 3D Voronoi/Delaunay

Voronoi cells are convex polyhedra.
Voronoi balls pass through 4 samples.

Delaunay
tetrahedra.


## Sliver tetrahedra



In 3D, some Voronoi vertices are not near medial axis ...

## Sliver tetrahedra


.... even when samples are arbitrarily dense.

Interior Voronoi balls

## Poles



Interior polar balls

Subset of Voronoi vertices, the poles, approximate medial axis.

Amenta \& Bern, 98 "Crust" papers

## Poles

For dense surface samples, Voronoi cells are:

- long and skinny,
- perpendicular to surface,
- with ends near the medial axis.



## Poles



Poles are Voronoi vertices at opposite ends.

To find: farthest Voronoi vertex from sample, farthest on opposite side.

## Sampling Requirement



ع-sample: distance from any surface point to nearest sample is at most small! constant $\varepsilon$ times distance to medial axis. Note: surface has to be smooth.

## Sampling Requirement

Intuition: dense sampling where curvature is high or near features.

## Large balls tangent

Any large ball (with respect to distance to medial axis) touching sample s has to be nearly tangent to the surface at $s$.

## Specifically

Given an $\varepsilon$-sample from a surface F:

Angle between normal to F at sample s and vector from $s$ to either pole $=0$ ( $\varepsilon$ )


## Results

Look for algorithms where....
Input: $\varepsilon$-sample from surface $G$
Output: PL-surface,

- near $G$, converges
- normals near $G$, converge
- PL manifold
- homeomorphic to $G$


## Formal Algorithms

Amenta and Bern, crust Amenta, Choi, Dey and Leekha, co-cone Boissonnat and Cazals, natural neighbor Amenta, Choi and Kolluri, power crust

## Co-cone

Estimate normals, choose candidate triangles with good normals at each vertex.

Extract manifold from candidates.

## Co-cone



Works well on clean data from a closed surface.

Amenta, Choi, Dey, Leekha 2000

## Co-cone extensions



Dey \& Giesen, undersampling errors.
Dey \& Goswami, hole-filling.

Dey, Giesen \& Hudson, divide and conquer for large data.

## Power Crust



Amenta, Choi and Kolluri, 01

Idea: Approximate object as union of balls, compute polygonal surface from balls.

## Power Crust



Start with all poles.

## Power Crust



Compute polygonal decomposition using power diagram.

## Power Crust



Label power diagram cells inside or outside object (skipping details).

Inside cells form polyhedral solid.

## Power Crust



Boundary of solid gives output surface.

Connect inner poles with adjacent power diagram cells for approximate medial axis.

## Example



Laser range data, power crust, simplified approximate medial axis.

## Medial axis approximation

Dey \& Zhao, 02 Voronoi diagram far from surface.


## Medial axis approximation

Medial axis of union of balls = lower dimensional parts of alpha shape + intersection with Voronoi diagram of union vertices.


Attali \& Montanvert, 97, A \& Kolluri, 01

## Distance function



> Giesen and John, 01,02

Distance from nearest sample.

## Distance function



Consdier uphill flow .... Idea: interior is part that flows to interior maxima.

## Distance function



Compute flow combinatorially using Delaunay/ Voronoi

Max and (some) saddle points.

## Distance function structure



Critical points where dual Delaunay and Voronoi faces intersect.

## Distance function structure



Critical points where dual Delaunay and Voronoi faces intersect.

## Not all pairs are critical



## Wrap

## Edelsbrunner - (95), Wrap, to appear....



Product!
Based on similar flow idea.


## Running time

All $O\left(n^{2}\right)$ in theory because of complexity of 3D Delaunay triangulation. Practically, Delaunay is bottleneck.

But not in practice?

## Delaunay complexity lower bound

Arrange points on two skew line segments $O\left(n^{2}\right)$ Delaunay triangulation


## Dimension of distribution

Dwyer, 91: Random points in sphere have Delaunay triangulation of size $O(n)$ in $R^{d}$.

$O(n)$

??

$O\left(n^{2}\right)$

## 2D surface in $R^{3}$

Cylinder
[Erickson01]

$O\left(n^{3 / 2}\right)$
$O(n \log n)$
$O(n)$

## Tangent spheres



Infinite number of tangent spheres touch an infinite number of surface points.


All but a finite number of tangent spheres touch a finite number of surface points


All tangent spheres touch a finite number of surface points

## Polyhedral ( $\mathrm{d}-1$ ) in $\mathrm{R}^{\mathrm{d}}$

Delauany triangulation of points on d-1 dimensional polyhedral surfaces:

\# simplices per sample:

4D-3
5D-20
6D - 100

## Generic (d-1) in $R^{d}$

Delauany of points on cubic polynomial


$$
\begin{aligned}
& \text { \# simplices } \\
& \text { per sample: } \\
& 4 D-5 \\
& 5 D-30 \\
& 6 D-130
\end{aligned}
$$

## Upper bound, polyhedral

[A,Attali\&Devillers] Number of
Delaunay simplicies for $n$ points nearly uniformly sampling all faces of a pdimensional polyhedral surface (not nec. connected or convex) in $R^{d}$ is:
$O\left(n^{(d+1-k) / p}\right), \quad k=\operatorname{ceiling}((d+1) /(p+1))$

## Next week...

...maybe we'll get a chance to improve this.

