

Quiz 2

First name: Last name: Seat #:

Instructions: No notes/books/gadgets/neighbors. Be mathematically precise.

1. Let's recall the **product construction**, as used to show that the DFA-acceptable languages are closed under **union**. Let $M_1 = (Q_1, \Sigma, \delta_1, q_1, F_1)$ be a DFA for L_1 and let $M_2 = (Q_2, \Sigma, \delta_2, q_2, F_2)$ be a DFA for L_2 . We construct a DFA $M = (Q, \Sigma, \delta, q_0, F)$ for $L_1 \cup L_2$ by asserting that $Q = \text{$, that $\delta((p, q), a) = \text{$, that $q_0 = \text{$, and, finally, that $F = \text{$. If $|Q_1| = N_1$, $|F_1| = n_1$, $|Q_2| = N_2$, and $|F_2| = n_2$, then $|F| = \text{$ (formula involving N_1, n_1, N_2, n_2).
2. Suppose we use the **subset construction** to convert a 7-state NFA $M = (Q, \Sigma, \delta, q_0, F)$ into a DFA $M' = (Q', \Sigma, \delta', q_0, F')$ for the same language. Then M' will have $|Q'| = \text{$ states (a number). If $|F| = 2$ then M' will have $|F'| = \text{$ final states (a number).
3. The string 00100 is the th binary string in lexicographic order (assume $0 < 1$).
4. **Darken** the **correct** box. No justification is required. If you're not sure, guess.

(a)	<input type="checkbox"/> True	<input type="checkbox"/> False	Every finite language is regular.
	<input type="checkbox"/> True	<input type="checkbox"/> False	If A and B are regular then $A - B$ is regular.
	<input type="checkbox"/> True	<input type="checkbox"/> False	If L^* is infinite the L is infinite.
	<input type="checkbox"/> True	<input type="checkbox"/> False	Regular expressions are strings.
	<input type="checkbox"/> True	<input type="checkbox"/> False	If $M = (Q, \Sigma, \delta, q_0, F)$ is an NFA then $x \in L(M)$ iff $\delta^*(q_0, x) \subseteq F$.
	<input type="checkbox"/> True	<input type="checkbox"/> False	If NFA $M = (Q, \Sigma, \delta, s, F)$ has $F = Q$ then $L(M) = \Sigma^*$.
	<input type="checkbox"/> True	<input type="checkbox"/> False	The pumping lemma can be used to show that languages are regular.
	<input type="checkbox"/> True	<input type="checkbox"/> False	Let $L = \{a^n b^n : n \geq 1\}$. Then L^* is regular.
	<input type="checkbox"/> True	<input type="checkbox"/> False	An n -state NFA M can be converted into an n^2 -state DFA for $L(M)$.
	<input type="checkbox"/> True	<input type="checkbox"/> False	A regular expression α can be converted into a $2 \alpha $ -state NFA for $L(\alpha)$.
	<input type="checkbox"/> True	<input type="checkbox"/> False	If $A \subseteq L \subseteq B$ and A and B are regular then L is regular.
	<input type="checkbox"/> True	<input type="checkbox"/> False	In the Myhill-Nerode theorem, we defined $x \sim_L x'$ by: $[x \in L \Leftrightarrow x' \in L]$.
5. Carefully state the **pumping lemma** for regular languages. Use any form of this lemma you like, but make sure it's nontrivial and all quantifiers are clear.
6. Draw an **NFA** for $L = \{aab, aba\}$. Your NFA must use **as few states and arrows** as possible.