

Quiz 3

First: LAST: Seat: Row:

1. A PDA $M = (Q, \Sigma, \Gamma, \delta, q_0, F)$ has $|Q| = 10$ states and $|\Sigma| = 2$ characters in the input alphabet and $|\Gamma| = 5$ characters in the tape alphabet. Then there are points in the domain of δ and points in the range of δ .
2. State the **Church-Turing thesis**:
3. Let A and B be languages. Define $A \leq_m B$ (A **many-one reduces** to B):
4. When we define the language $A_{\text{TM}} = \{\langle M, w \rangle : \text{TM } M \text{ accepts } w\}$, what is the purpose of the **angle brackets** (the $\langle \rangle$ symbols) that surround M, w ?
5. **Darken** the **correct** box. No justification is required. If you're not sure, guess.
 - (a) True False If L is recursive then so is its complement, \bar{L} .
 - (b) True False If L^* is recursive then L is recursive.
 - (c) True False If L is context free then a queue automata (QA) can decide it.
 - (d) True False The r.e. languages are closed under complement.
 - (e) True False $L = \{\langle M \rangle : L(M) \neq \emptyset\}$ is Turing-acceptable (r.e.)
 - (f) True False $L = \{a^n b^n : n \geq 1\}$ is co-r.e.
 - (g) True False If $\Pi \leq_m L$ and Π is undecidable then L is undecidable.
 - (h) True False To show that L is not r.e., it suffice to show that $A_{\text{TM}} \leq_m L$.
 - (i) True False To show that L is not r.e., it suffice to show that $\overline{A_{\text{TM}}} \leq_m L$.
 - (j) True False A language L is either r.e. or co-r.e..
 - (k) True False The Turing-acceptable languages are closed under intersection.
 - (l) True False The Turing-acceptable languages are closed under set difference.