

## Quiz 9 Solutions

For this quiz I want you to prove that

$$A = \{\langle M, k \rangle : M \text{ is a TM that accepts at least one string of length } k\}$$

is undecidable. Do this with a reduction involving  $A_{\text{TM}}$  or  $\overline{A_{\text{TM}}}$ . Make your proof succinct, legible, and logical. Write exclusively in grammatical English sentences.

*Setup.* Since  $A$  is r.e., we will show that it is undecidable by showing that  $A_{\text{TM}} \leq_m A$ . To do this, we must construct a Turing-computable function that maps a string  $\langle M, w \rangle$  to a string  $\langle M', k \rangle$  such that TM  $M$  accepts  $w$  if and only if TM  $M'$  accepts some string of length  $k$ .

*Construction.* Given  $\langle M, w \rangle$  the reduction returns  $\langle M', k \rangle$  where  $k \geq 0$  is an arbitrary fixed value, say  $k = 0$ , and TM  $M'$  is the following machine:

Machine  $M'$ , on input  $x$ :  
Run  $M$  on  $w$   
If  $M$  accepts then accept  
If  $M$  rejects then reject

*Analysis.* If  $M$  accepts  $w$  then we will have that  $L(M') = \Sigma^*$ , so  $M'$  will accept a string of length  $k$  (as it accepts all strings of all lengths). On the other hand, if  $M$  does not accept  $w$  then  $L(M') = \emptyset$  so  $M'$  will not accept any string of length  $k$  (as it accepts no string of any length). Finally, the function that computes  $\langle M', k \rangle$  from  $\langle M, w \rangle$  is clearly computable.