# Practice Midterm Exam

**Instructions:** Relax. Smile. Be happy. Then **think** about each question for a few minutes before writing down a **brief**, **correct** answer!

Bon courage!

— Phil Rogaway

Name:

E-mail:

On problem	you got	out of
1		40
2		35
3		25
Σ		100

#### 1 Short Answer

## [40 points]

(1) Draw a DFA that accepts  $L = \{x \in \{a, b\}^* : x \text{ starts and ends with different characters}\}.$ 

(2) Using the procedure we have seen in class, convert the following NFA into a DFA that accepts the same language.



(3) You are given a regular expression  $\alpha$ . Describe a decision procedure (algorithm) that determines if  $L(\alpha)$  contains an odd-length string.

(4) Using the construction given in class and in your text, convert the regular expression  $\alpha = (01 \cup 1)^*$  into an NFA for the same language.

(5) Complete the definition (as given in class or your book): A context free grammar is a 4-tuple  $G = (V, \Sigma, R, S)$  where<sup>1</sup>

(6) Use closure properties to show that  $L = \{0^i 1^j 2^j : i, j \ge 0\}$  is not regular.

<sup>&</sup>lt;sup>1</sup>Don't just tell me what  $V, \Sigma, R, S$  are called; tell me what they *are*, mathematically.

(7) Let REG be the language of all (fully parenthesized) regular expressions over the alphabet  $\{0, 1\}$ .

Thus sample strings in REG are:

$$\varepsilon$$
  
1  
((0 \circ 1)  $\cup$  1)

Prove that REG is CF by giving a CFG for it.

(8) With REG as defined above, prove that REG is not regular by using the pumping lemma.

### 2 Justified True or False

Put an **X** through the **correct** box. Where it says "**Explain**" provide a **brief** (but convincing) justification. No credit will be given to correct answers that lack a proper justification. Where appropriate, **make your justification a counter-example**. Throughout, we use L and R to denote languages.

1.	If $\overline{L}$ is finite then L is regular.	True	False
	Explain:		
2.	Every regular language can be accepted by an NFA that has ex	actly 1,000,0	00 states
	Explain:	True	False

3.	Let $R$ be regular and let $L$ is a subset of $R$ . Then $L$ is regular		
	Explain:	True	False

[35 points]

4. If L is finite then L\* is infinite.Explain:

5. Suppose L has the following property: for some number N, every string  $s \in L$  having length at least N can be partitioned into s = xyz such that  $|y| \ge 1$  and  $xy^i z \in L$  for all  $i \ge 0$ . Then L is regular. **True False Explain**:

6. Every CFL L can be generated by a CFG G in which every rule  $A \rightarrow \alpha$  satisfies  $|\alpha| \leq 2$ . Explain:

7. For L regular, let d(L) be the number of states in a smallest DFA for L, and let n(L) be the number of states in a smallest NFA for L. Then for any regular language L,  $n(L) \leq (d(L))^2$ . Explain:

False

True

#### 3 A Closure Property of Regular Languages [25 points]

If L is a language over an alphabet  $\Sigma$  let

Two-Less(L) = { $y \in \Sigma^*$ : for some string x having  $|x| \le 2$ ,  $xy \in L$  }.

**Part A.** (A warm-up, just to make sure you understand the definition.) Is one a subset of the other: L and Two-Less(L)? If so, which is a subset of which?

**Part B.** (A warm-up, just to make sure you understand the definition.) Let  $P = \{1^2, 1^3, 1^5, 1^7, 1^{11}, 1^{13}, \dots\}$  be the set of prime numbers, encoded in unary. What's the shortest string in 1<sup>\*</sup> which is not in *Two-Less*(*P*)?

**Part C.** (Now the main problem. This has nothing to do with Parts A or B.) Prove that if L is regular, than so is Two-Less(L). (Describe any construction you use both in clear English and by a formal definition.)